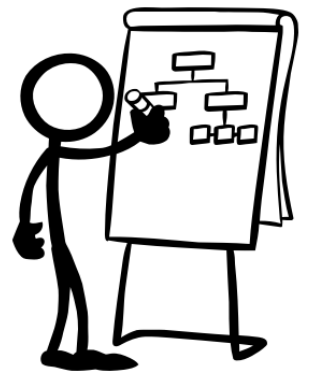


The Sweep Line Paradigm

Computational Geometry – Recitation 2



Agenda

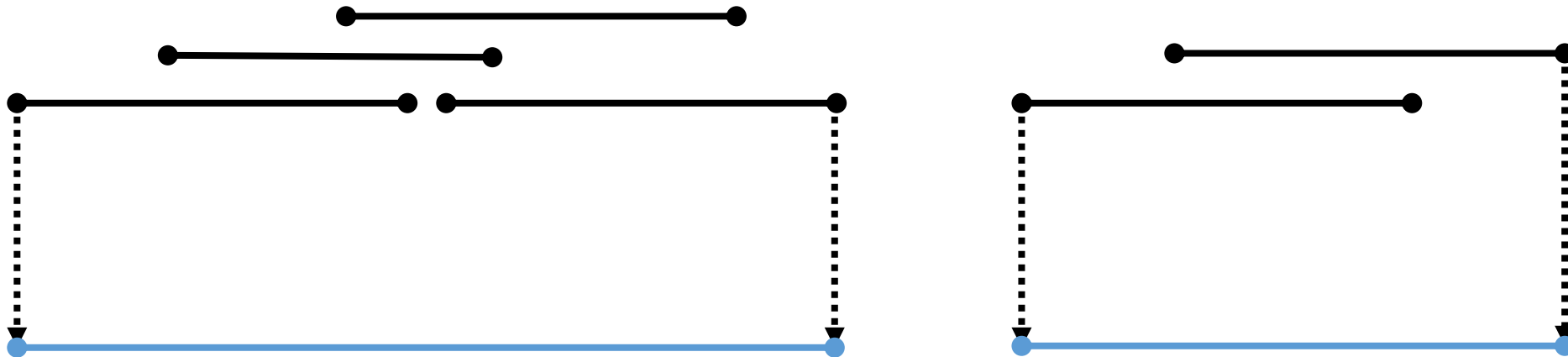
- Toy examples
- Line segment intersection
- Applications
 - Area of union of rectangles
 - Minimal distance pair



Sweeping: Example #1

Sweeping: Example #1

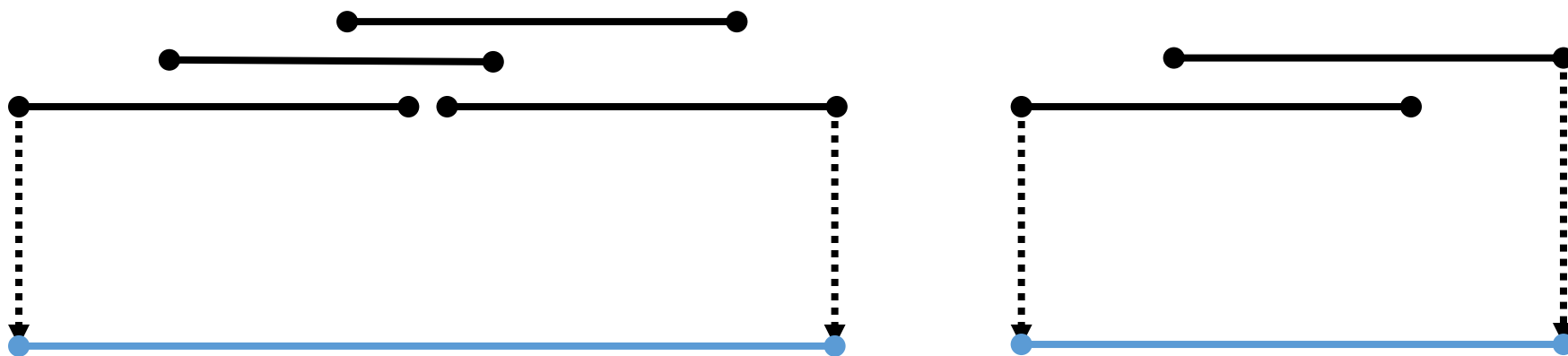
- Given a set of 1D segments, what is the union of them all?



- Solution: Sort all the points, and count the number of 'active' segments.

Sweeping: Example #1

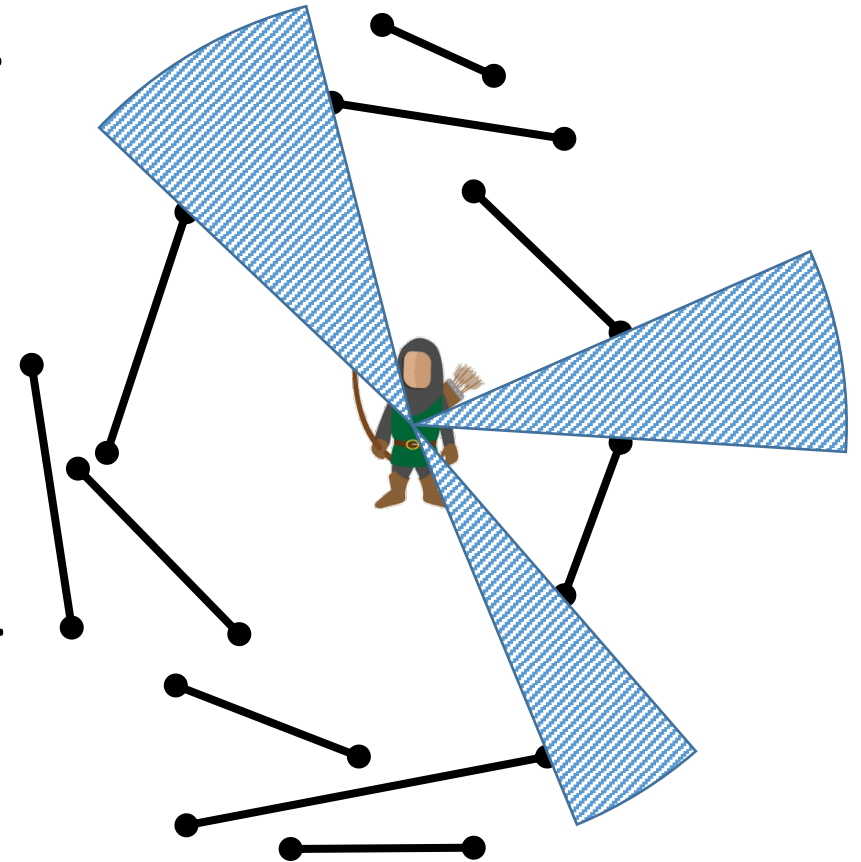
- We have traversed a discrete set of **Events**, in a certain **Order**, while maintaining some **Status** of the algorithm.
- **Events [What data was processed]**: start of segment, end of segment.
- **Order [In what order we traverse the events]**: From left to right
- **Status [Additional information maintained]**: number of active segments.
- Complexity: $O(n \log n)$



Sweeping: Example #2

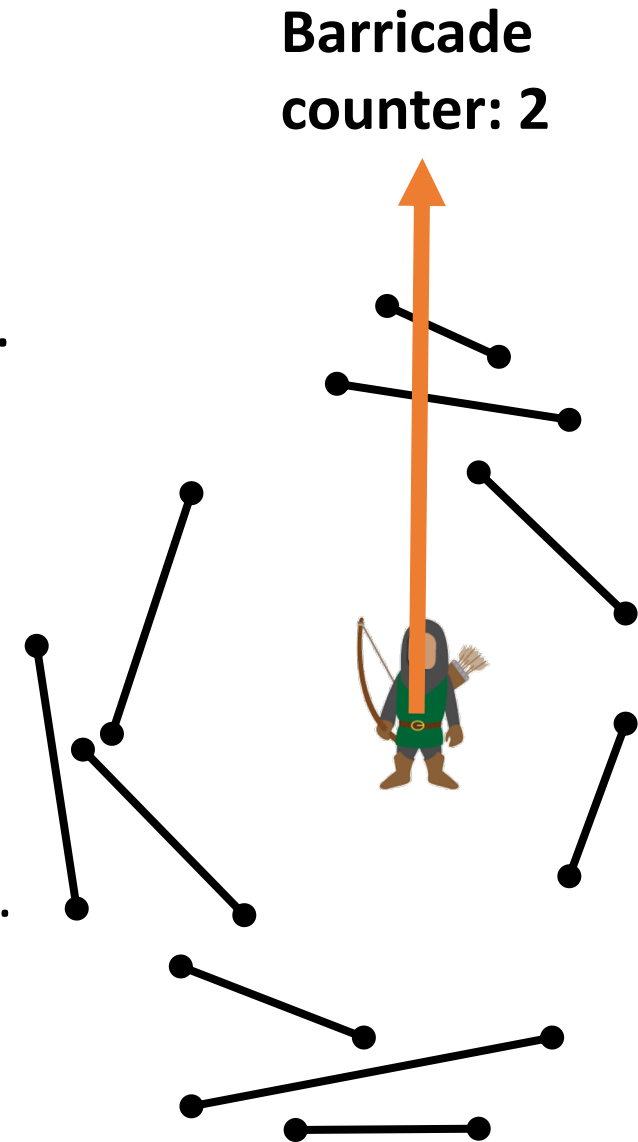
Sweeping: Example #2

- An archer is surrounded by a set of barricades. What are his lines of sight?
- **Order:** Scan the segments by angle.
- **Status:** Number of 'active' barricades.
 - Init in $O(n)$.
- **Events:**
 - Start of a segment: increase number of barricades.
 - End of a segment: decrease number of barricades.
- Report angles with 0 barricades.



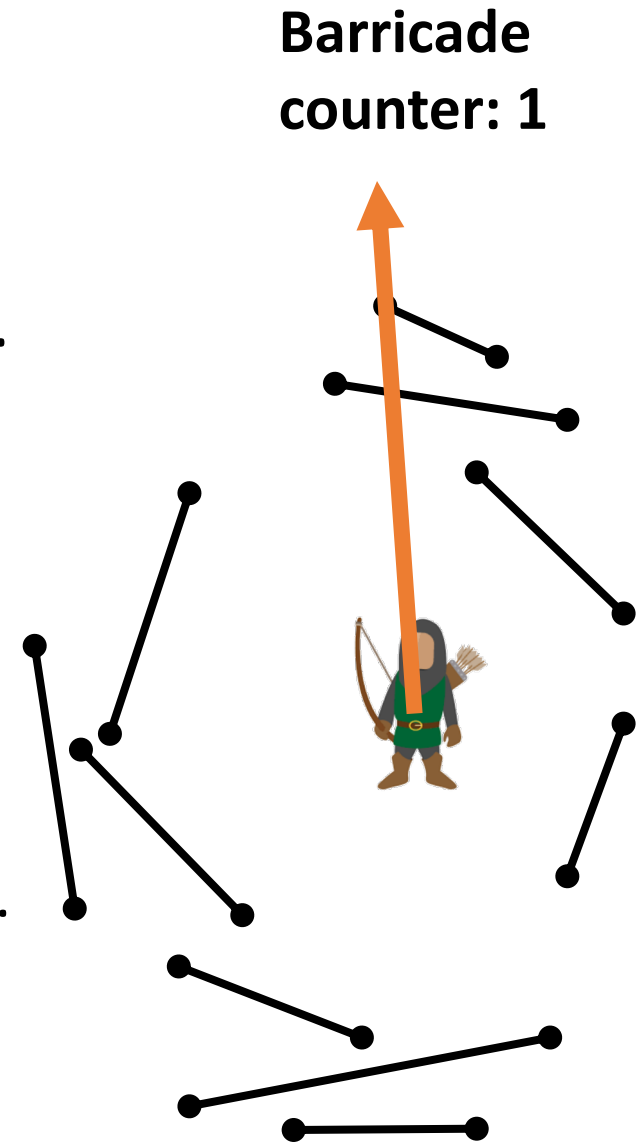
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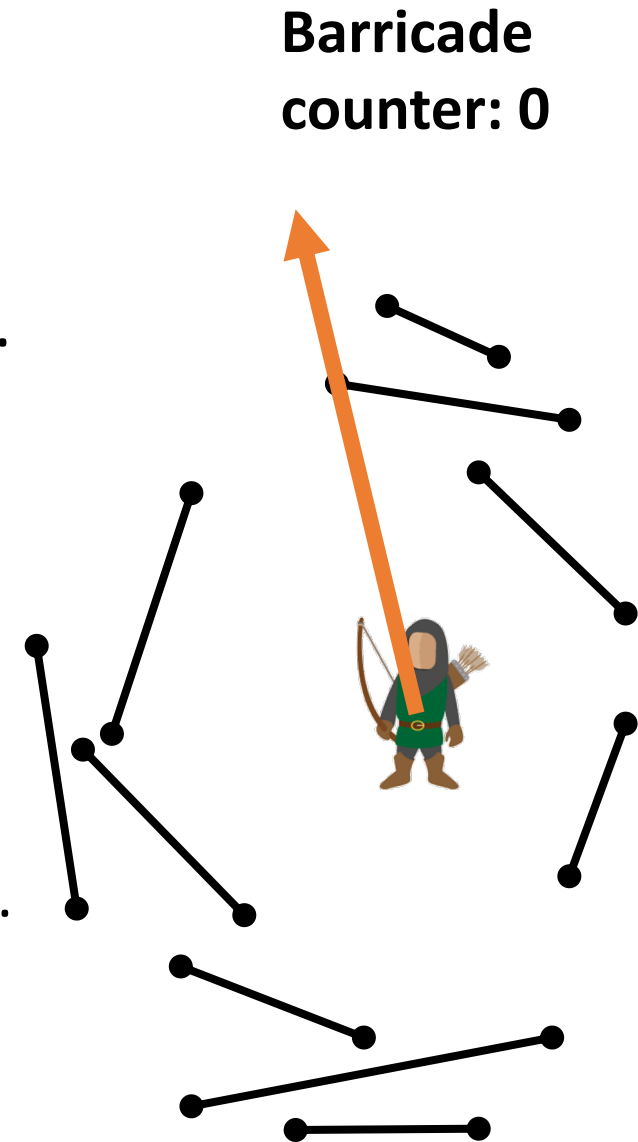
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Sweeping: Example #2

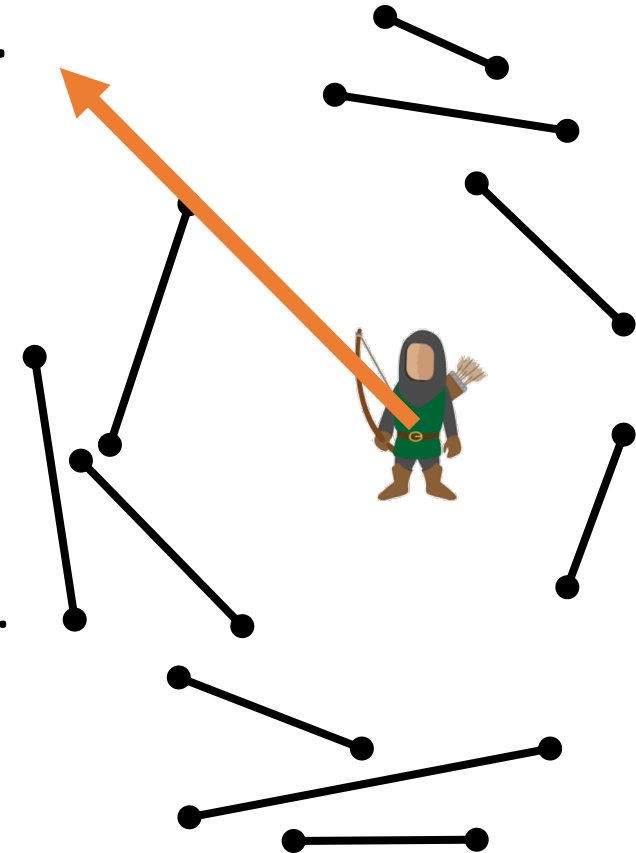
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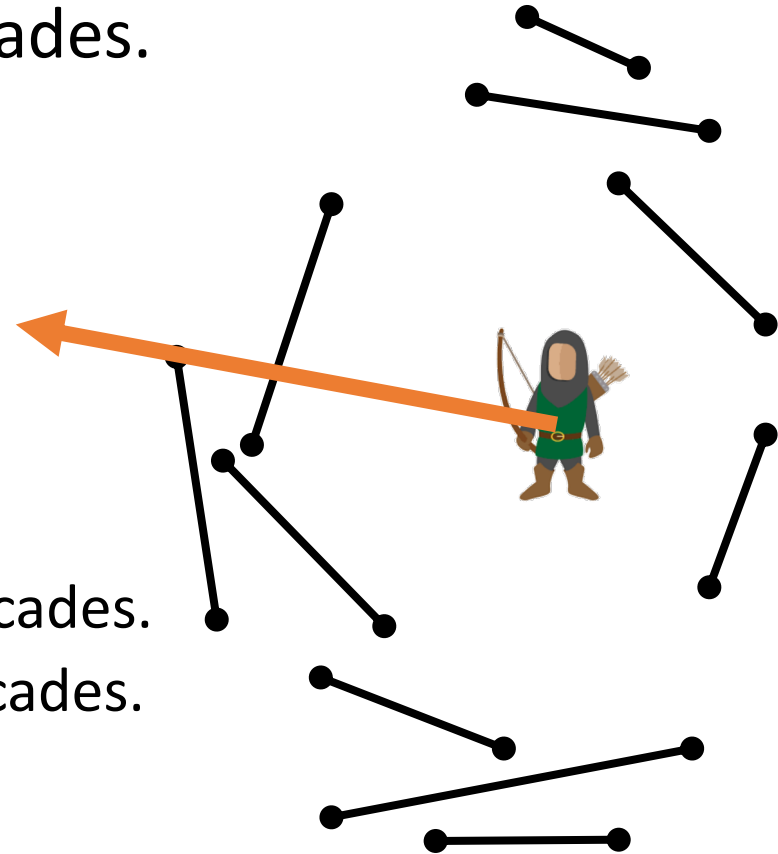
Barricade
counter: 1



Sweeping: Example #2

- An archer is surrounded by a set of barricades. What are his lines of sight?
- **Order:** Scan the segments by angle.
- **Status:** Number of 'active' barricades.
 - Init in $O(n)$.
- **Events:**
 - Start of a segment: increase number of barricades.
 - End of a segment: decrease number of barricades.
- Report angles with 0 barricades.
- Complexity: $O(n \log n)$

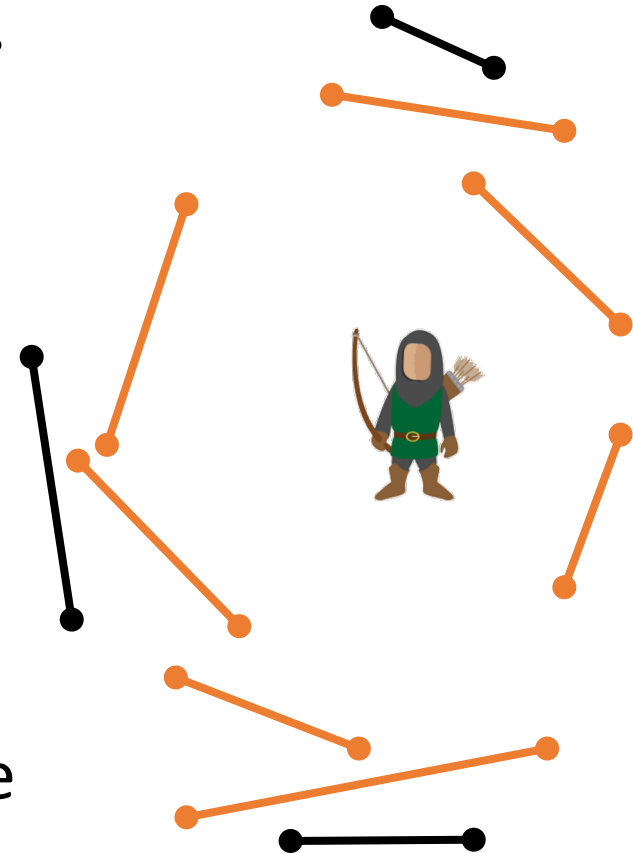
Barricade
counter: 2



Sweeping: Example #3

Sweeping: Example #3

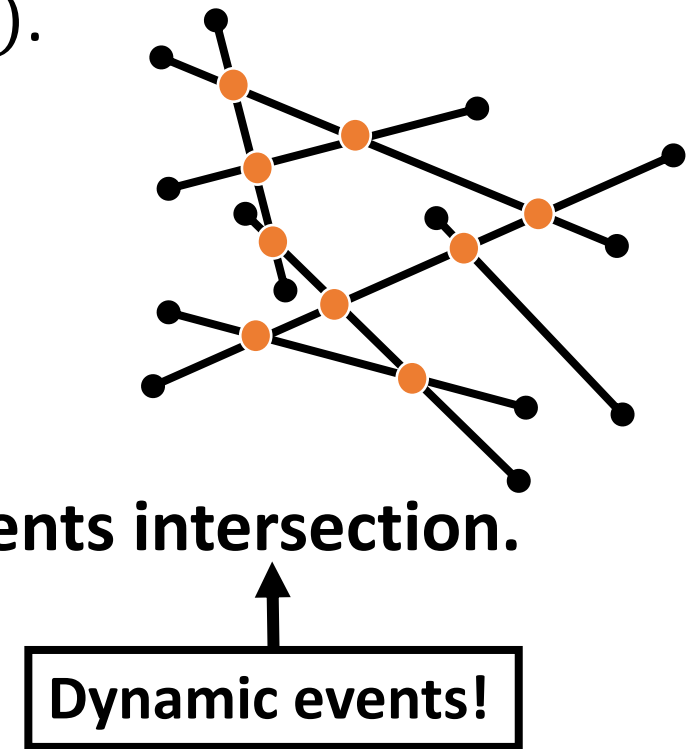
- An archer is surrounded by a set of barricades. Which barricades are visible to him?
- **Order:** Scan the segments by angle.
- **Status:** Set of active barricades, sorted by the distance from the archer.
- **Events:**
 - Start of a segment: Add segment to the status DS.
 - End of a segment: Remove segment from the status DS.
- Report all segments which was closest at some point.
- Complexity: $O(n \log n)$



Segment Intersection

Segment Intersection

- Given a set of n segments, report all intersection points.
- Naïve algorithm: Check all segment pairs, $O(n^2)$.
- Sweep line algorithm:
- **Order:** scan from left to right.
- **Status:** segments intersecting the sweep line.
(Ordered by intersection point).
- **Events:** Segment start, Segment end and **Segments intersection.**
- **Check intersection only between adjacent segments in the status DS.**



Handle event: *None*

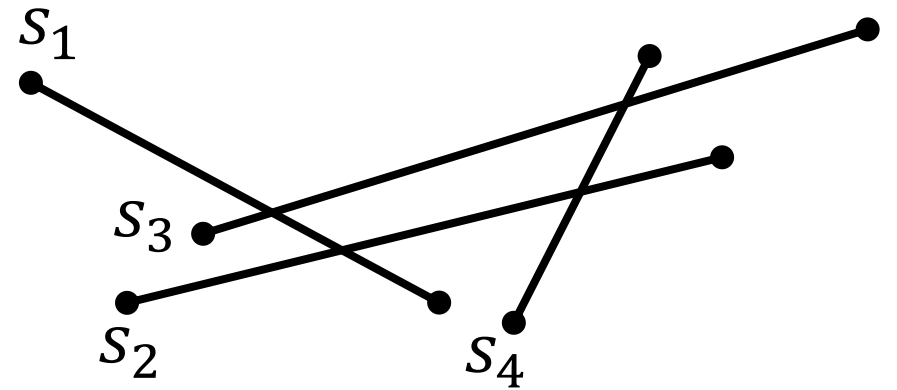
Events

$Start(S_1)$
$Start(S_2)$
$Start(S_3)$
$End(S_1)$
$Start(S_4)$
$End(S_4)$
$End(S_2)$
$End(S_3)$

Status

\emptyset

Sweep line

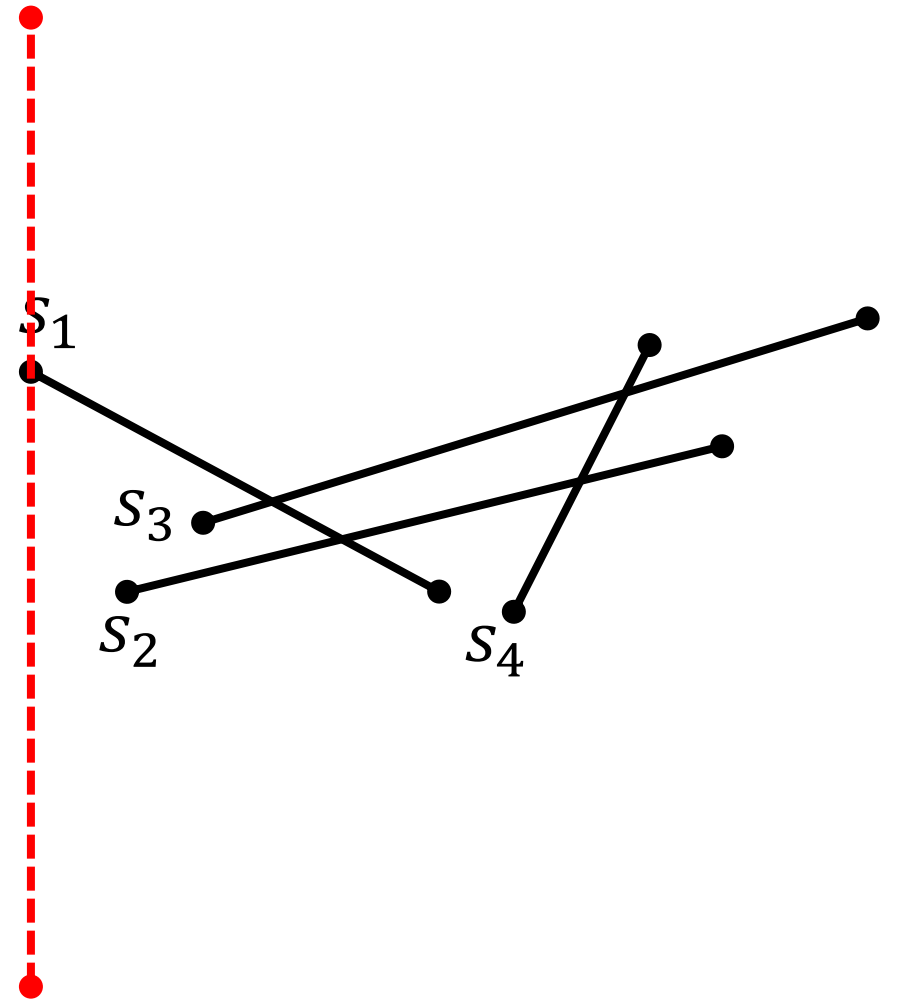
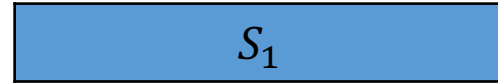


Handle event: $Start(S_1)$

Events

$Start(S_2)$
$Start(S_3)$
$End(S_1)$
$Start(S_4)$
$End(S_4)$
$End(S_2)$
$End(S_3)$

Status



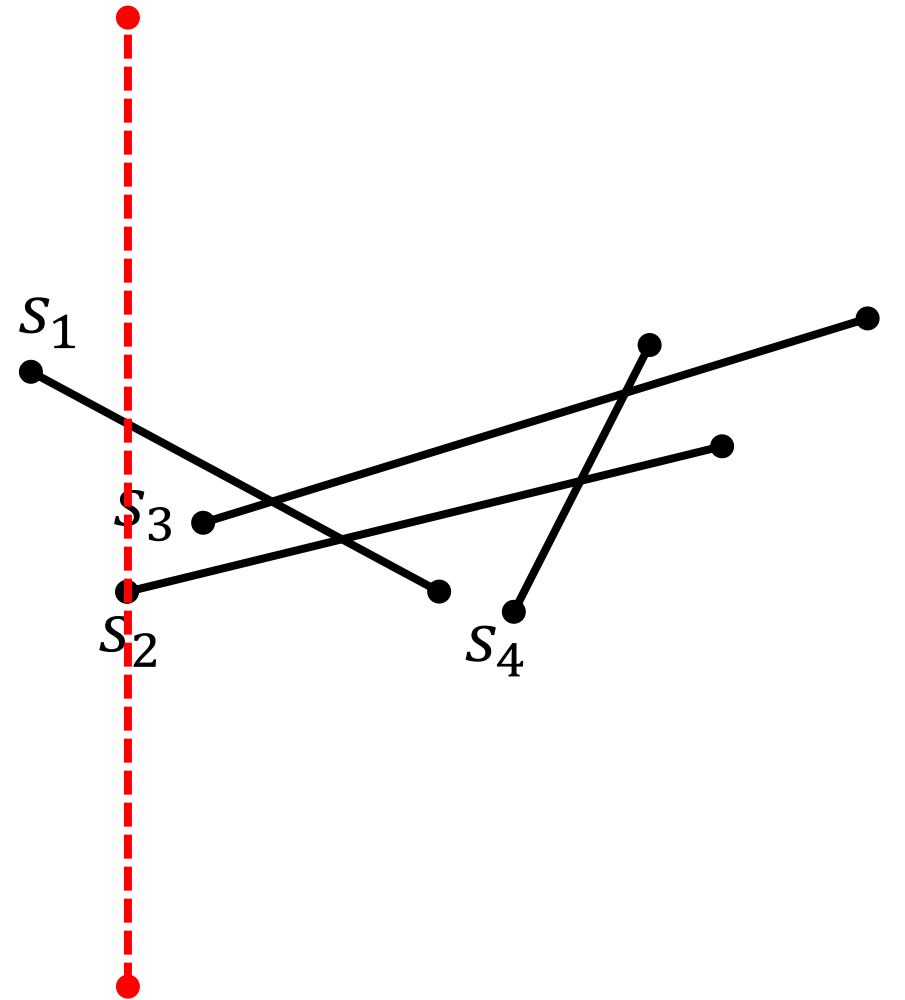
Handle event: $Start(S_2)$

Events

$Start(S_3)$
$Intersection(S_1, S_2)$
$End(S_1)$
$Start(S_4)$
$End(S_4)$
$End(S_2)$
$End(S_3)$

Status

S_1
S_2



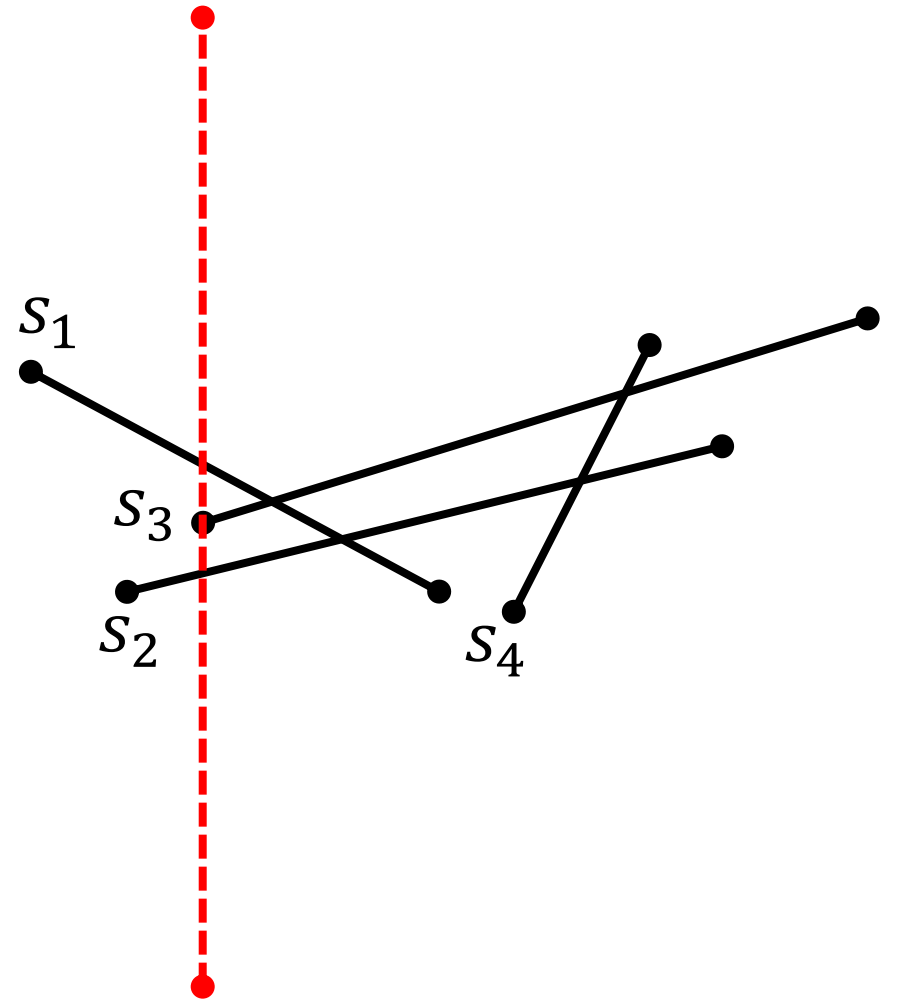
Handle event: $Start(S_3)$

Events

$Intersection(S_1, S_3)$
$Intersection(S_1, S_2)$
$End(S_1)$
$Start(S_4)$
$End(S_4)$
$End(S_2)$
$End(S_3)$

Status

S_1
S_3
S_2



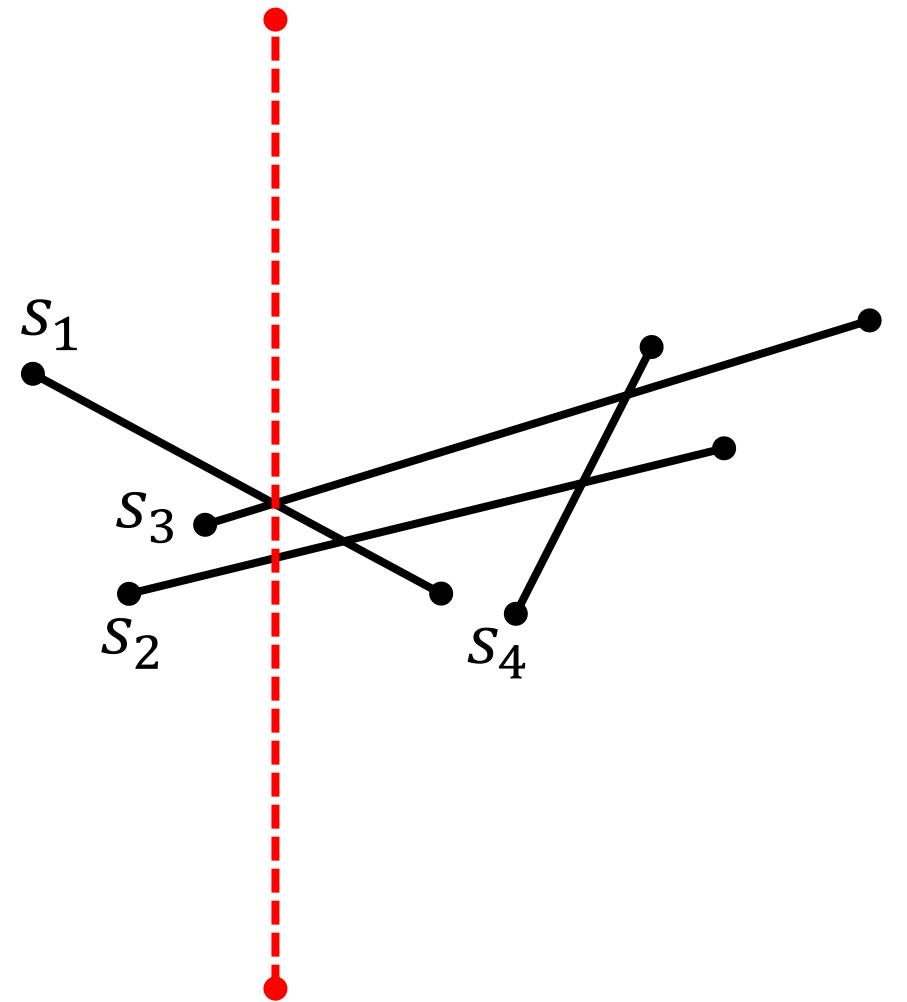
Handle event: $Intersection(S_1, S_3)$

Events

$Intersection(S_1, S_2)$
$End(S_1)$
$Start(S_4)$
$End(S_4)$
$End(S_2)$
$End(S_3)$

Status

S_3
S_1
S_2



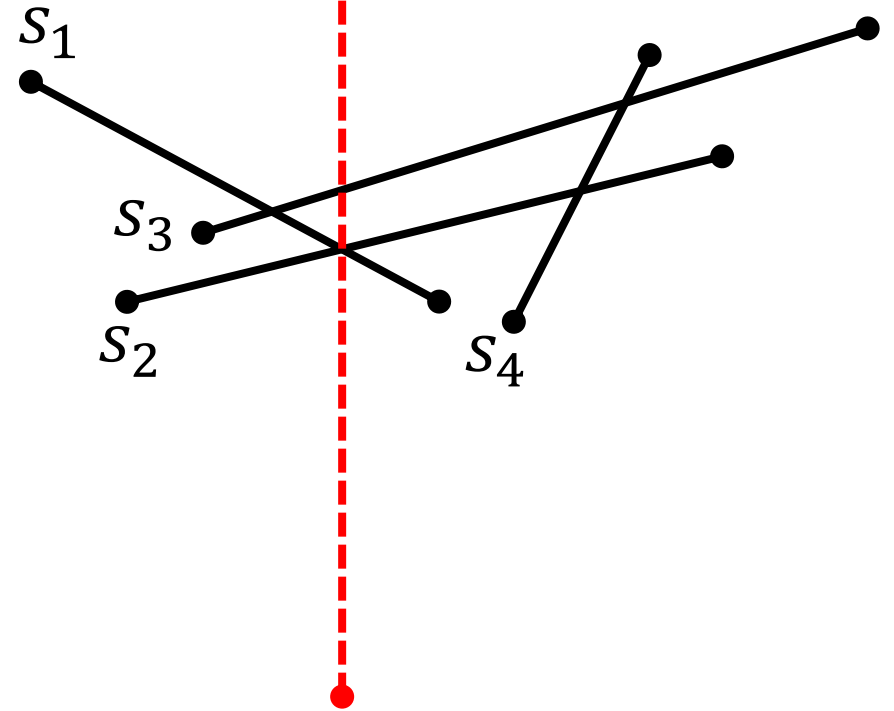
Handle event: $Intersection(S_1, S_2)$

Events

$End(S_1)$
$Start(S_4)$
$End(S_4)$
$End(S_2)$
$End(S_3)$

Status

S_3
S_2
S_1



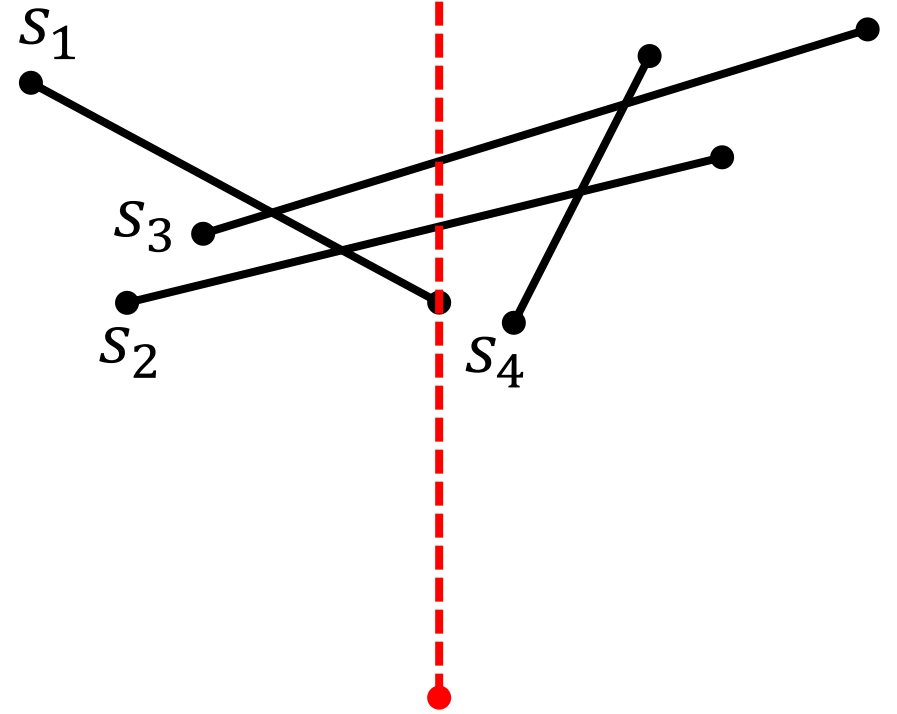
Handle event: $End(S_1)$

Events

$Start(S_4)$
$End(S_4)$
$End(S_2)$
$End(S_3)$

Status

S_3
S_2



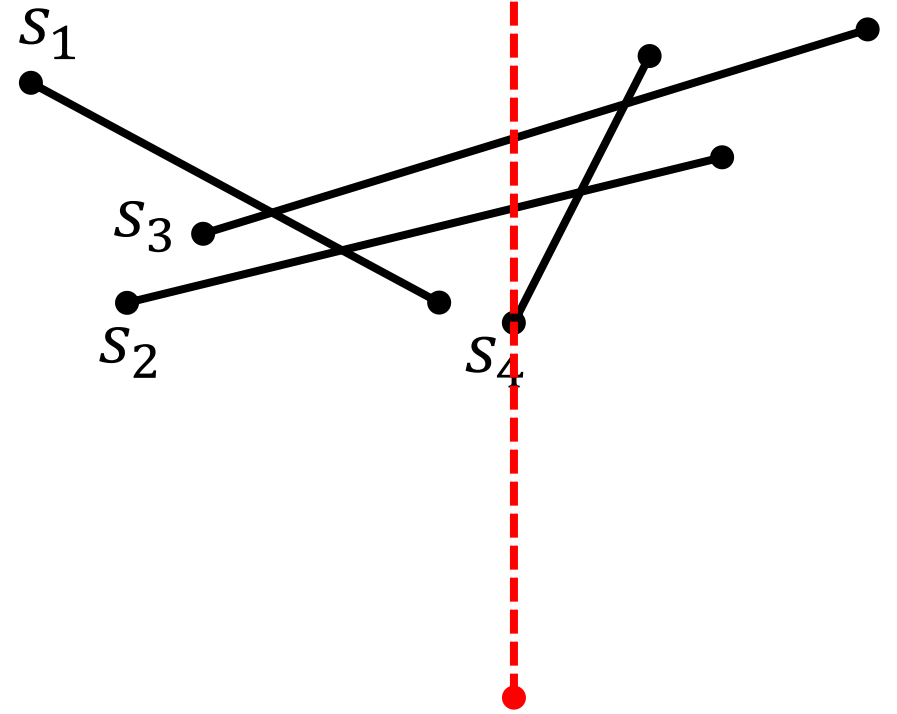
Handle event: $Start(S_4)$

Events

$Intersection(S_2, S_4)$
$End(S_4)$
$End(S_2)$
$End(S_3)$

Status

S_3
S_2
S_4



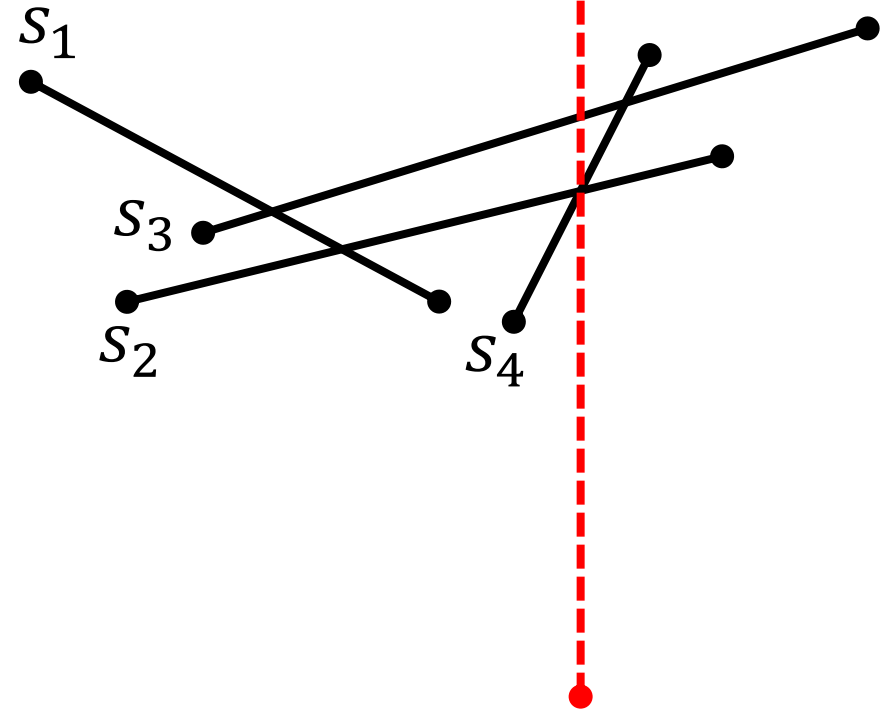
Handle event: $Intersection(S_2, S_4)$

Events

$Intersection(S_3, S_4)$
$End(S_4)$
$End(S_2)$
$End(S_3)$

Status

S_3
S_4
S_2



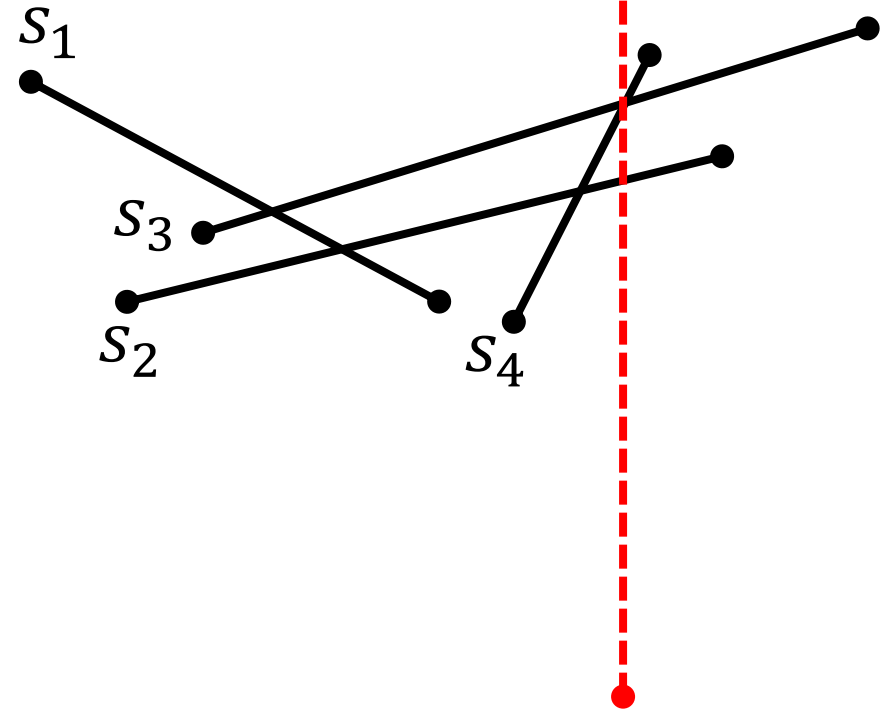
Handle event: $Intersection(S_3, S_4)$

Events

$End(S_4)$
$End(S_2)$
$End(S_3)$

Status

S_4
S_3
S_2



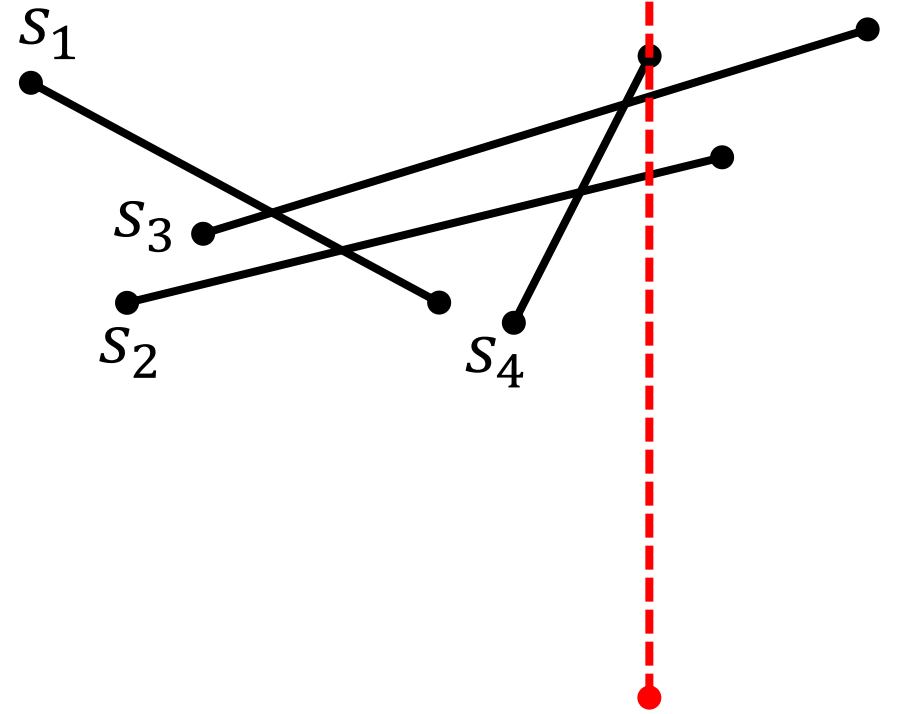
Handle event: $End(S_4)$

Events

$End(S_2)$
$End(S_3)$

Status

S_3
S_2



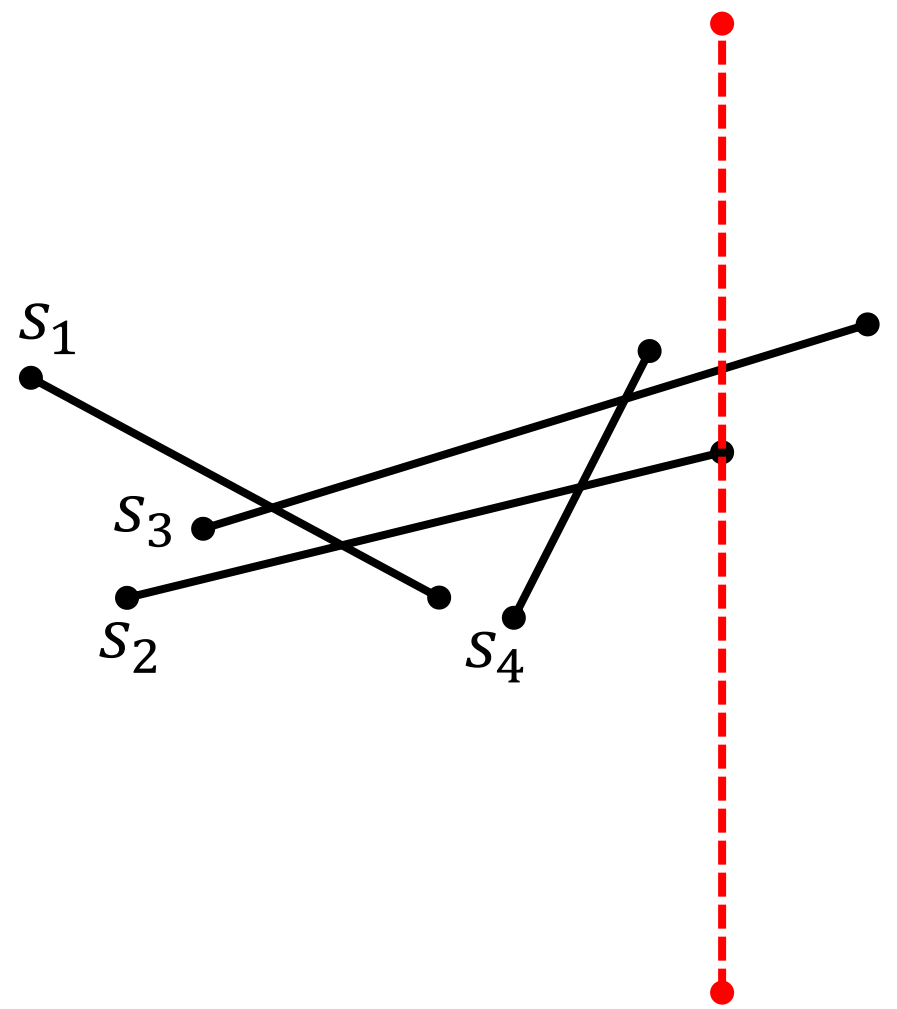
Handle event: $End(S_2)$

Events

$End(S_3)$

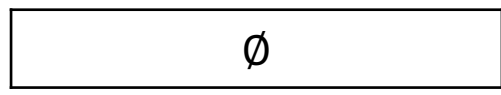
Status

S_3

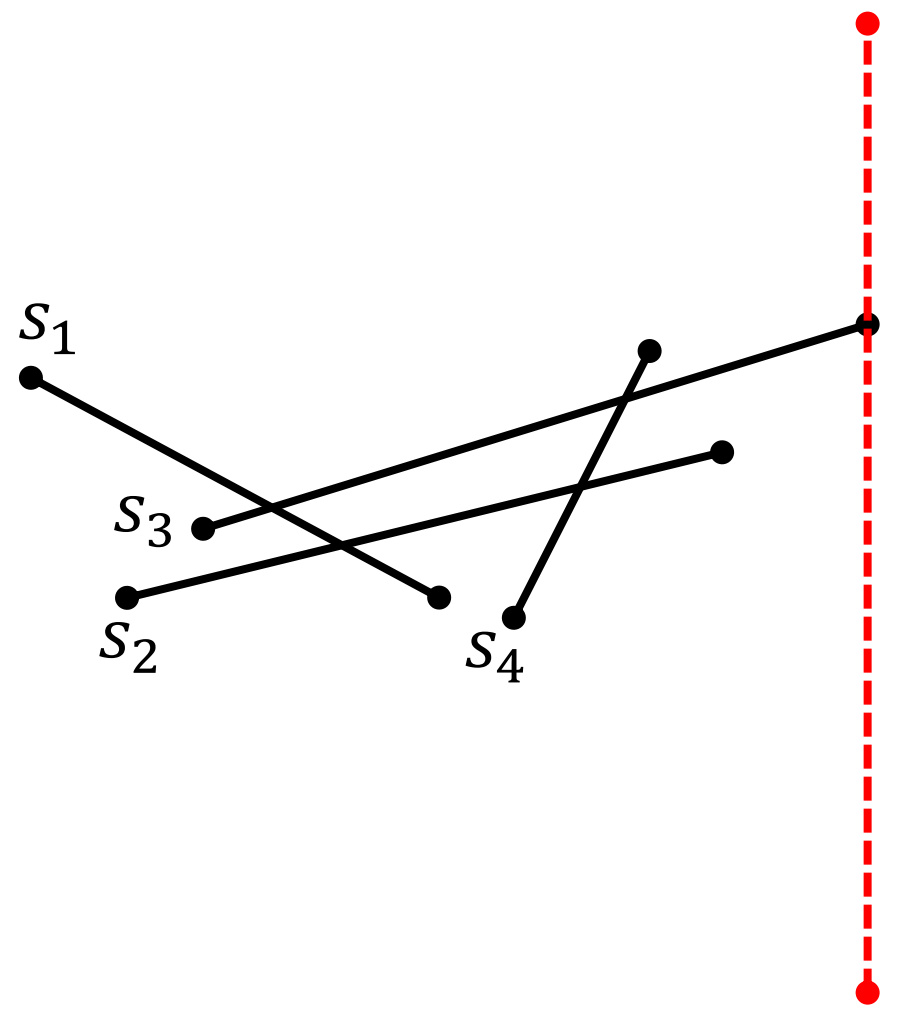
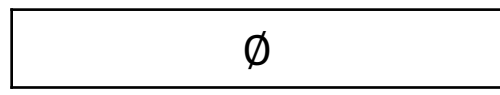


Handle event: $End(S_3)$

Events

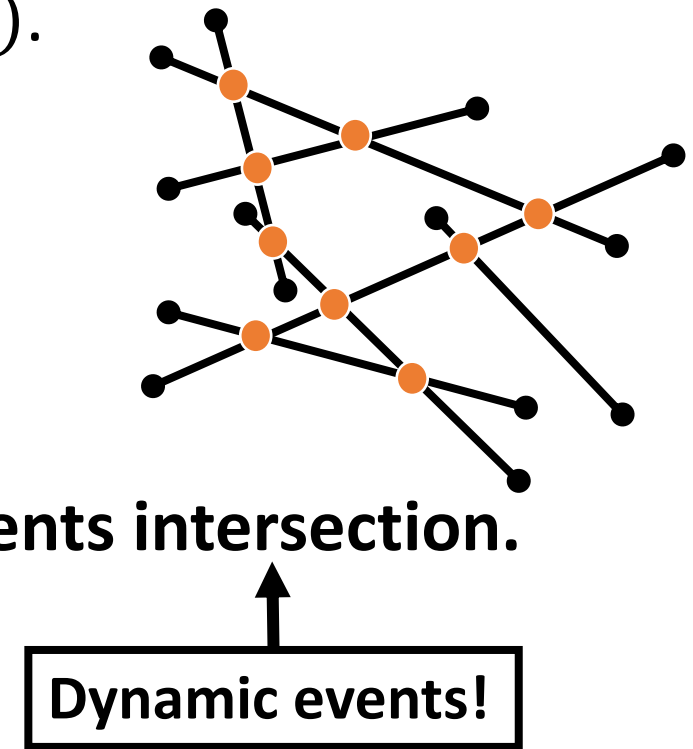


Status



Segment Intersection

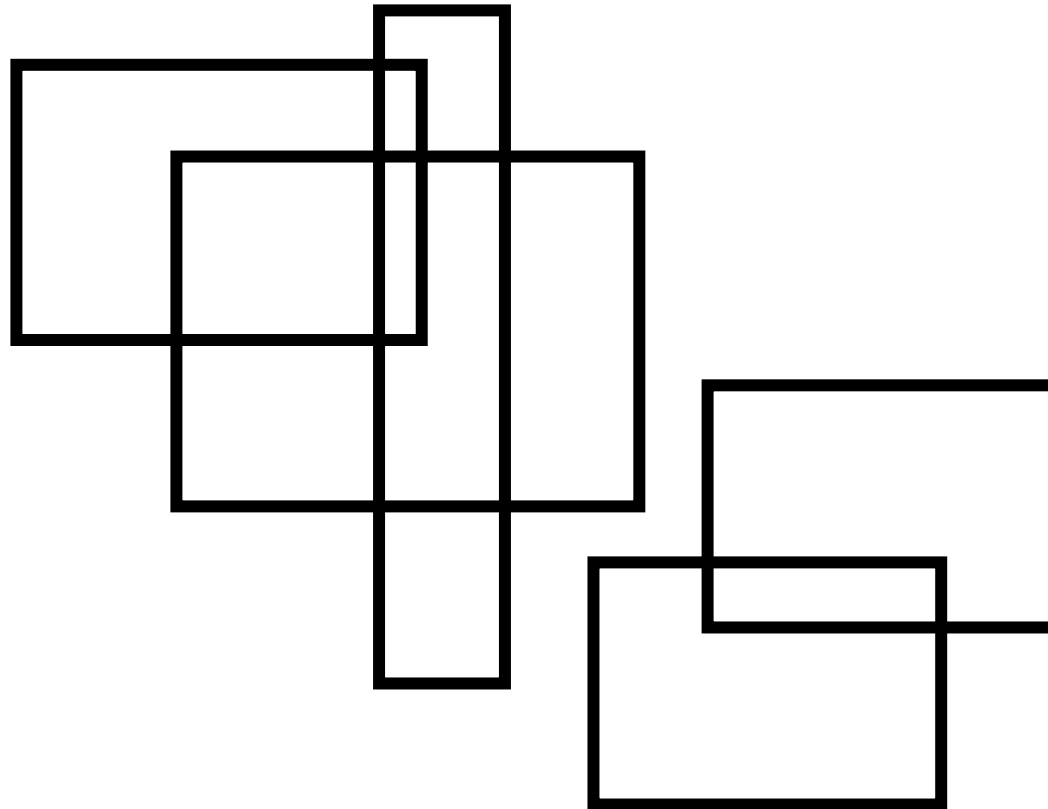
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 - **Events:** Segment start, Segment end and **Segments intersection.**
 - **Check intersection only between adjacent segments in the status DS.**
- Complexity: $O(n \log n)$



Area of union of Rectangles

Area of union of Rectangles

- What is the total area covered by a set of rectangles?



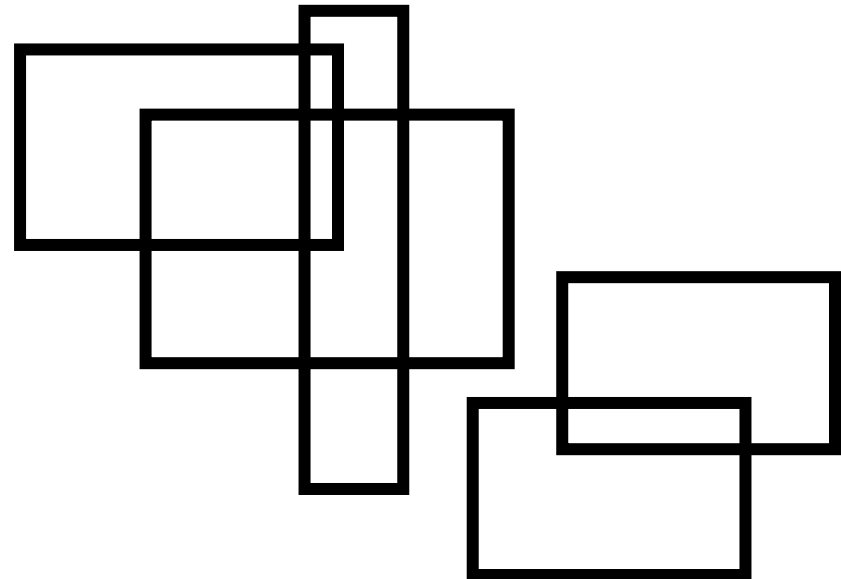
Area of union of Rectangles

- What is the total area covered by a set of rectangles?

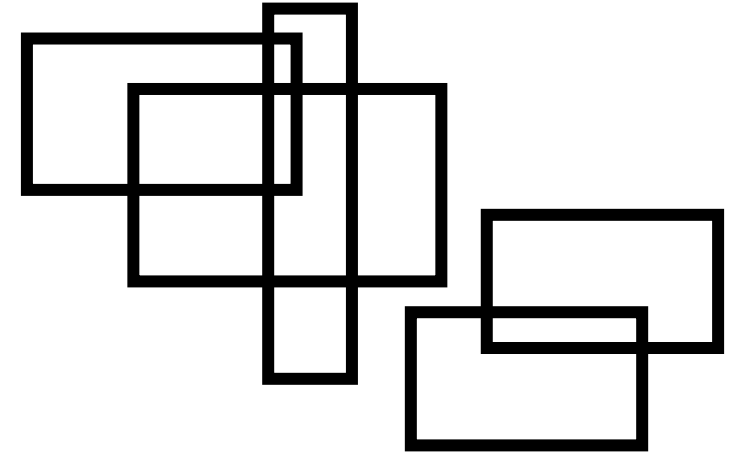


Area of union of Rectangles

- What is the total area covered by a set of rectangles?
- **Order:** left to right
- **Events:** begin and end of a rectangle
- **Status:** active rectangles



Area of union of Rectangles

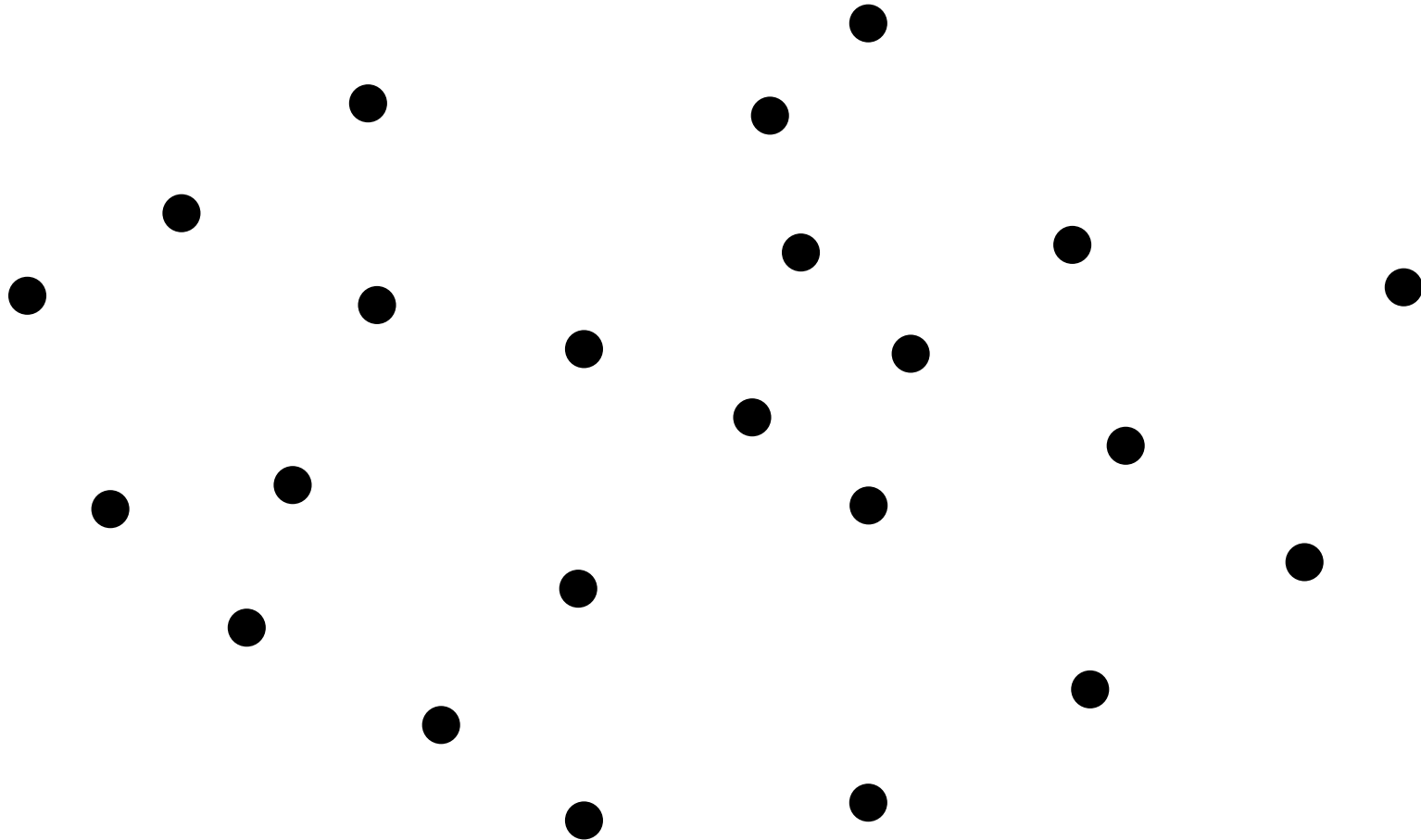


- **Status:** active rectangles
- How do we maintain the active rectangle set?
- More importantly, how do we find the total length covered by the active rectangle?
- Naïve implementation: Recalculate the union each time (using example #1).
Complexity: $O(n^2)$.
- Better implementation: Use augmented BST (classic DS exercise).
Complexity: $O(n \log n)$.

Minimal Distance Pair

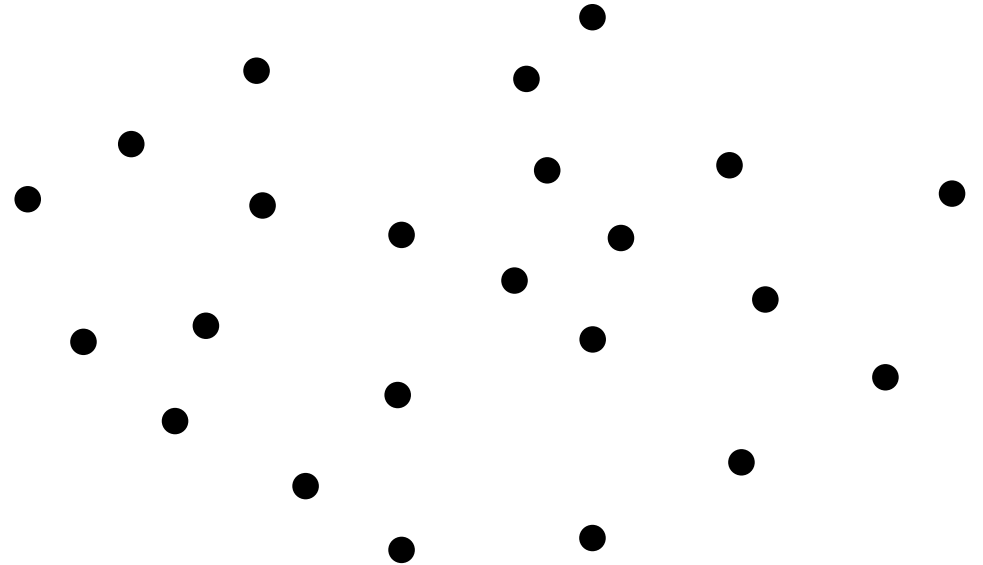
Minimal Distance Pair

- Problem: Find the closest pair of points.

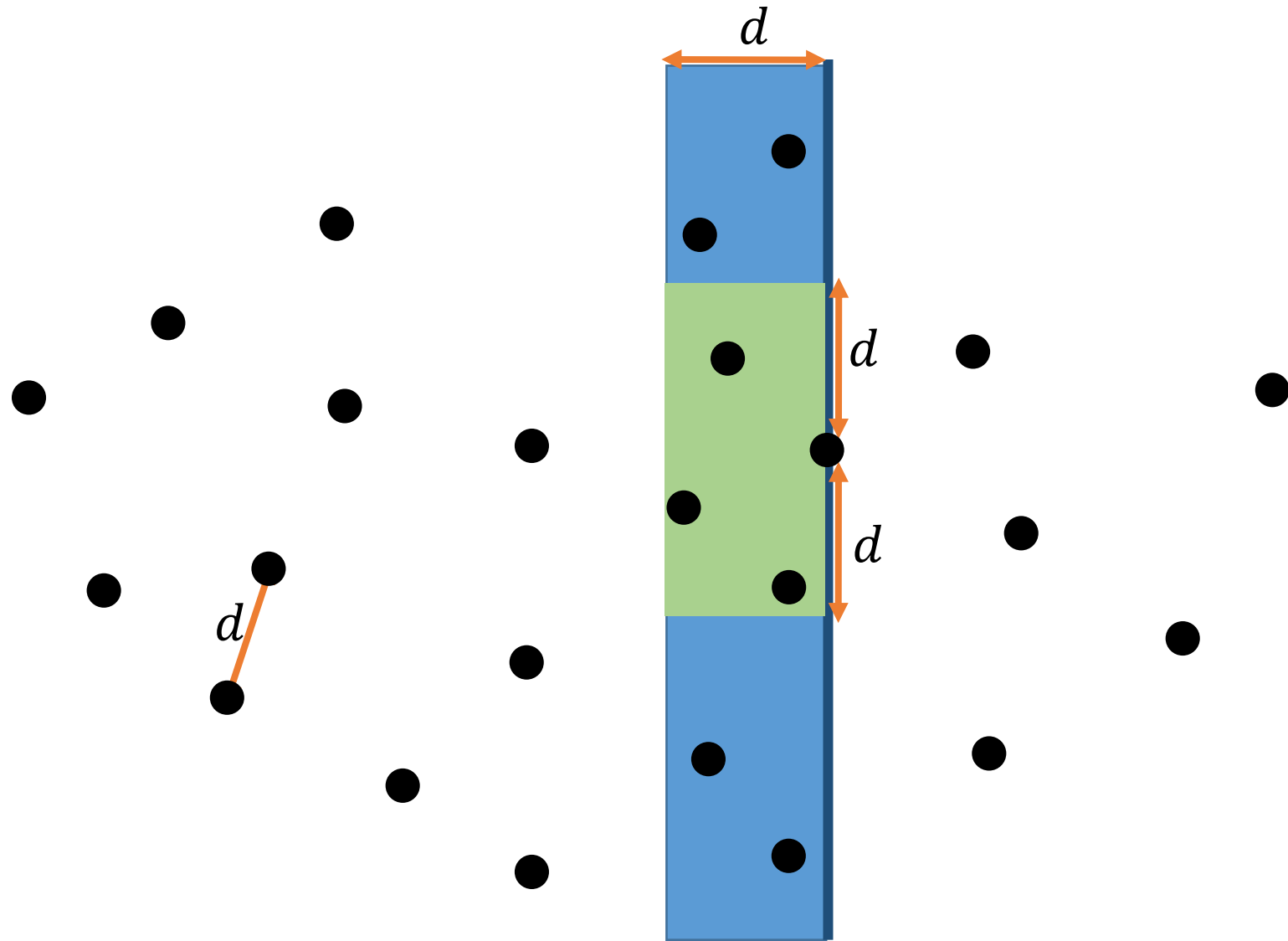


Minimal Distance Pair

- Problem: Find the closest pair of points.
- Naïve algorithm: Check all pairs, $O(n^2)$
- Sweeping idea:
- **Events:** All the points
- **Order:** left to right
- **Status:** minimal distance seen so far, d .
 And two BSTs of all the points in a strip of width d .
 one sorted by the y coordinate,
 and another sorted by the x coordinate.



Minimal Distance Pair



Minimal Distance Pair

- Handle event:
- Compare the distance with the relevant points.
 - Using the sorted by y tree.
- Update d if needed.
- Remove from both trees the points that now are not part of the strip.
 - Using the sorted by x tree.